

Chapter 3: Discrete Random Variables

3.1 The Notion of a Random Variable

3.1

Sample Space:
 Coins

Michael		0	1	2
$\frac{1}{4}$	0	(0,0)	(0,1)	(0,2)
$\frac{1}{2}$	1	(1,0)	(1,1)	(1,2)
$\frac{1}{4}$	2	(2,0)	(2,1)	(2,2)
		$\frac{1}{4}$	$\frac{1}{2}$	$\frac{1}{4}$

Probabilities

		0	1	2
0		$\frac{1}{16}$	$\frac{1}{8}$	$\frac{1}{16}$
1		$\frac{1}{8}$	$\frac{1}{4}$	$\frac{1}{8}$
2		$\frac{1}{16}$	$\frac{1}{8}$	$\frac{1}{16}$

Mapping $\mathcal{S} \rightarrow \mathcal{X}$

		0	1	2
0		0	1	2
1		1	1	2
2		2	2	2

$P[X=0] = P[(0,0)] = \frac{1}{16}$
 $P[X=1] = P[\{(1,0), (1,1), (0,1)\}] = \frac{1}{2}$
 $P[X=2] = 3 \times \frac{1}{16} + 2 \times \frac{1}{8} = \frac{7}{16}$

3.2

(a) $S = \{1, 2, 3, 4, 5, 6\}$ $p_1 = p_2 = p_3 = p_4 = p_5 = p_6 = \frac{1}{6}$
 where $p_j = P[\{j\}]$

(b)

S		$\sum X$	
1	→	0	
2	→	1	
3	→	1	
4	→	2	
5	→	2	
6	→	3	

(c)

$P[X=0] = p_1 = \frac{1}{6}$
 $P[X=1] = p_2 + p_3 = \frac{2}{6}$
 $P[X=2] = p_4 + p_5 = \frac{2}{6}$
 $P[X=3] = p_6 = \frac{1}{6}$

(d) $P[X=0] = p_1 + p_2 = \frac{2}{6}$ $P[Y=1] = p_3 + p_4 = \frac{2}{6}$ $P[Y=2] = p_5 + p_6 = \frac{2}{6}$

(e) $X=0$ corresponds to $\{1\}$
 $Y=0$ corresponds to $\{1, 2\}$

3.3

(a) $A = \{(x, y) : x^2 + y^2 = r^2\}$ where $r = \text{radius of circle}$

Outcomes from A occur uniformly along the circle.
 $\text{sgn}(xy) = 0$ at the dots

(c)

$P[X=-1] = P[2\text{nd} + 3\text{rd} \text{ Quad}] = \frac{1}{2}$
 $P[X=0] = P[\{(r, 0), (0, r), (-r, 0), (0, -r)\}] = 0$
 $P[X=1] = P[1\text{st} + 3\text{rd} \text{ Quad}] = \frac{1}{2}$

2nd & 3rd Quadrant
 4 dots
 1st & 3rd Quadrant

-1 0 +1

3.4

a) $S = \{0000, 0001, \dots, 1111\}$
 $p_{0000} = p_{0001} = \dots = p_{1111} = \frac{1}{16}$

b)

S	0000	0001	0010	...	1111
	↓	↓	↓		↓
S_x	0	1	2	...	15

c) $p_0 = p_1 = p_2 = \dots = p_{15} = \frac{1}{16}$

d)

$P[0b_1b_2b_3]$	$= \frac{1}{4} \cdot \frac{1}{8} = \frac{1}{32}$	all $b_1b_2b_3$
$P[1b_1b_2b_3]$	$= \frac{3}{4} \cdot \frac{1}{8} = \frac{3}{32}$	all $b_1b_2b_3$

$p'_0 = p'_1 = \dots = p'_7 = \frac{1}{32}$

$p'_8 = p'_9 = \dots = p'_{15} = \frac{3}{32}$

3.5 Let $A_i =$ Transmitter #1 sends a signal at time slot i
 $B_i =$ " #2 " "

A signal gets through if $A_i \cdot B_i^c \cup A_i^c \cdot B_i$ occurs

Each experiment has 4 outcomes

(a) Experiment i

	A_i	A_i^c
B_i	$\frac{1}{4}$	$\frac{1}{4}$
B_i^c	$\frac{1}{4}$	$\frac{1}{4}$

Sample Space consists of a Cartesian product of the outcomes of the basic experiment

$S = (s_1, s_2, \dots)$ where s_i is an outcome from basic experiment

(b) $X(s) = n$
 if n is the first occurrence of $A_i \cdot B_i^c \cup A_i^c \cdot B_i$ in s_1, s_2, \dots

(c) $P[A_i \cdot B_i^c \cup A_i^c \cdot B_i] = P[A_i \cdot B_i^c] + P[A_i^c \cdot B_i] = \frac{1}{2} = P[\text{success}]$
 $P[X=k] = P[(k-1) \text{ failures, 1 success}] = \left(\frac{1}{2}\right)^k$

3.6 $\mathcal{A} = \{000, 111, 010, 101, 001, 110, 100, 011\}$

$X(s):$ \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow \downarrow
 2 2 3 3 4 4 4 4

$P[X=2] = P[\{000, 111\}] = \frac{1}{2}$
 $P[X=3] = P[\{010, 101\}] = \frac{1}{4}$
 $P[X=4] = P[\{001, 110, 100, 011\}] = \frac{1}{4}$

3.7) Draw 2 bills without replacement.

		2nd bill				
		1_1	1_2	...	1_9	50
1st bill	1_1	x				
	1_2		x		2	
	⋮			⋮		
	1_9					x
	50				51	x

x not allowed
 all other outcomes
 have probability $\frac{1}{9(10)} = \frac{1}{90}$

$$P[X=2] = \frac{9 \cdot 8}{90} = \frac{8}{10} = \frac{4}{5}$$

$$P[X=51] = \frac{9 \cdot 2}{90} = \frac{2}{10} = \frac{1}{5}$$

3.8) Draw 2 bills with replacement.

		2nd bill				
		1_1	1_2	...	1_9	50
1st bill	1_1					
	⋮			2		
	1_9					
	50				51	100

all outcomes have
 probability $\frac{1}{10(10)} = \frac{1}{100}$

$$P[X=2] = \frac{81}{100}$$

$$P[X=51] = \frac{18}{100}$$

$$P[X=100] = \frac{1}{100}$$

3.9) (a) Let m be number of tails $0 \leq m \leq n$
 then number of heads is $n-m$ and the difference is
 $Y = n-m-m = n-2m \quad 0 \leq m \leq n$
 $\therefore S_Y = \{-n, -n+2, \dots, n-2, n\}$

(b) $P[Y=0] = P[n=2m] = P\left[m = \frac{n}{2}\right]$ for n even.

$P[Y=k] = P[n-2m=k] = P\left[m = \frac{n-k}{2}\right]$ for $n-k$ even

3.10

Let $S = \{b_1, b_2, \dots, b_{2^m}\}$ be the sequence of
 m -bit passwords as covered by the hacker.
 The target system picks a password at random from S .
 $X(S)$ is the index of the selected password.

$S_X = \{1, 2, \dots, 2^m\}$ where the value of X is
 selected at random from S_X .

$P[i] = \frac{1}{2^m} \quad i \in S_X$.

3.2 Discrete Random Variables And Probability Mass Function

3.11

(a)

the max function shifts probability mass to higher values of k

(b) If Carlos uses a biased coin:

	Carlos				
	0	1	2		
Midel	0	0	1	2	$\frac{1}{4}$
	1	1	1	2	$\frac{1}{2}$
	2	2	2	2	$\frac{1}{4}$
	$\frac{1}{16}$	$\frac{6}{16}$	$\frac{9}{16}$		

$$P[X'=0] = \frac{1}{4} \cdot \frac{1}{16} = \frac{1}{64}$$

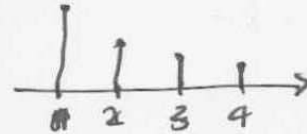
$$P[X'=1] = \frac{1}{16} \cdot \frac{1}{2} + \frac{6}{16} \cdot \frac{1}{2} + \frac{6}{16} \cdot \frac{1}{4} = \frac{20}{64}$$

$$P[X'=2] = \frac{43}{64}$$

3.12

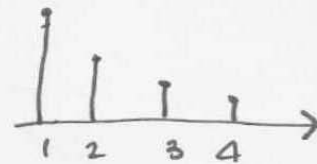
$$(a) \quad 1 = p_1 + p_2 + p_3 + p_4 = p_1 \left(1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4}\right) = \frac{25}{12} p_1 \quad p_1 = \frac{12}{25}$$

$$p_1 = \frac{12}{25} \quad p_2 = \frac{6}{25} \quad p_3 = \frac{4}{25} \quad p_4 = \frac{3}{25}$$



$$(b) \quad 1 = p_1 \left(1 + \frac{1}{2} + \frac{1}{4} + \frac{1}{8}\right) = \frac{15}{8} p_1$$

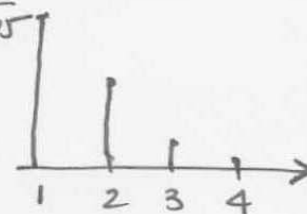
$$p_1 = \frac{8}{15} \quad p_2 = \frac{4}{15} \quad p_3 = \frac{2}{15} \quad p_4 = \frac{1}{15}$$



$$(c) \quad 1 = p_1 \left(1 + \frac{1}{2} + \frac{1}{8} + \frac{1}{64}\right) = \frac{105}{64} p_1$$

$$p_1 = \frac{64}{105} \quad p_2 = \frac{32}{105} \quad p_3 = \frac{8}{105} \quad p_4 = \frac{1}{105}$$

pmf decays more steeply w/
each example



(d) $1 = p_1 \sum_{i=1}^{\infty} \frac{1}{i}$ does not converge so this pmf
does not extend to $\{1, 2, \dots\}$

$$1 = p_1 \sum_{i=0}^{\infty} \left(\frac{1}{2}\right)^i = p_1 \frac{1}{1 - \frac{1}{2}} \Rightarrow p_1 = \frac{1}{2}$$

this extends to the geometric pmf.

$$1 = p_1 \left(1 + \left(\frac{1}{2}\right) + \left(\frac{1}{2}\right)^{1+2} + \left(\frac{1}{2}\right)^{1+2+3} + \dots\right)$$

$$= p_1 \sum_{j=1}^{\infty} \left(\frac{1}{2}\right)^{j(j+1)/2}$$

this is a subseries of
the geometric series
so it converges.

3.13

(a) $1 = \sum_{k=1}^{\infty} \frac{c}{k^2} = c \sum_{k=1}^{\infty} \frac{1}{k^2} = \frac{\pi^2}{6}$ is a special case of the zeta function
 $= 1.6449 \Rightarrow c = 0.608$

The sum of the first 100 terms gives $1.6349 \Rightarrow c \approx 0.611$

(b) $P[X > 4] = 1 - P[X \leq 3] = 1 - c \left[1 + \frac{1}{4} + \frac{1}{9} \right]$
 $= 0.1675$

(c) $P[6 \leq X \leq 8] = c \left[\frac{1}{36} + \frac{1}{49} + \frac{1}{64} \right] = 0.390$

3.14

$P[X \geq 8] = \sum_{k=8}^{15} p_k = \frac{8}{16} = \frac{1}{2}$

$P[Y \geq 8] = \sum_{k=8}^{15} p'_k = 8 \cdot \frac{3}{32} = \frac{24}{32} = \frac{3}{4}$

3.15

		Terminal 2	
		P	1-P
Terminal 1	$\frac{1}{2}$	$\frac{1}{2}P$	$\frac{1}{2}q$
$\bar{1}$	$\frac{1}{2}$	$\frac{1}{2}P$	$\frac{1}{2}q$

$P_{\text{success}} = \frac{1}{2}q + \frac{1}{2}P = \frac{1}{2}$ same

\therefore The pmf of X is unchanged.

$$P[\text{Terminal 2 transmitted} | \text{success}] = \frac{P[\text{success and Terminal 2 transmitted}]}{P[\text{success}]}$$

$$= \frac{\frac{1}{2}P}{\frac{1}{2}} = P$$

This suggests that terminal 2 should always transmit (at the expense of terminal 1).

3.16) from problem 3.7b:

(a) $P[X > 2] = 1 - P[X=2] = \frac{1}{5}$
 $P[X > 50] = P[X=51] = \frac{1}{5}$

(b) $P[X > 2] = 1 - P[X=2] = \frac{19}{100}$
 $P[X > 50] = P[X=51] + P[X=100] = \frac{19}{100}$

3.17

(a) $Y = 0 + 2 = 2$ with prob. $\frac{4}{10}$
 $Y = -1 + 2 = 1$ " $\frac{3}{10}$
 $Y = -2 + 2 = 0$ " $\frac{2}{10}$
 $Y = -3 + 2 = -1$ " $\frac{1}{10}$

(b) $P[Y=2] = \frac{4}{10}$

(c) $P[Y > 0] = P[Y=2] + P[Y=1] = \frac{4}{10} + \frac{3}{10} = \frac{7}{10}$

3.18) Let X be number of transmissions until ~~not~~ success.

$$P[X \leq k] = \sum_{j=1}^k \left(\frac{1}{2}\right)^j = \frac{1}{2} \sum_{j=0}^{k-1} \left(\frac{1}{2}\right)^j = \frac{1}{2} \frac{1 - \left(\frac{1}{2}\right)^k}{\frac{1}{2}} = 1 - \left(\frac{1}{2}\right)^k$$

$1 - \left(\frac{1}{2}\right)^k = 0.99$ $\left(\frac{1}{2}\right)^k = 0.01$

$k = \frac{\ln 100}{\ln 2} = 6.64 \approx 7$

start sending refresh messages
 7x 10 seconds before expiry time

3.19 $P[\text{decoding error}] = P[3 \text{ or more bit errors}]$

$$= \binom{5}{3} p^3 (1-p)^2 + \binom{5}{4} p^4 (1-p) + \binom{5}{5} p^5$$

$$= \frac{5!}{2!3!} 10^{-3} (0.9)^2 + \frac{5!}{4!1!} 10^{-4} (0.9) + 10^{-5}$$

$$= (0.81)(10)(10^{-3}) + (0.9)(5)10^{-4} + 10^{-5}$$

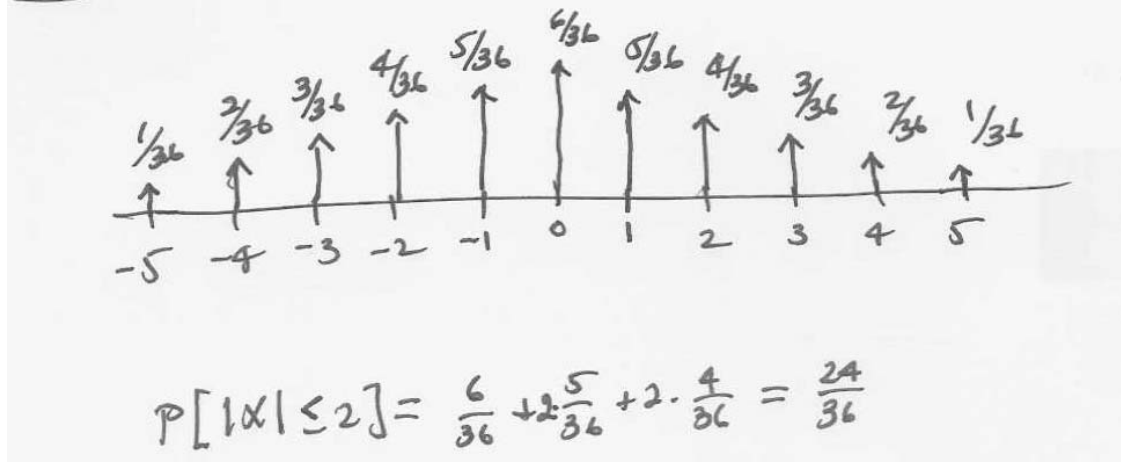
$$= 0.00856$$

which is an order of magnitude less than without coding.

3.20

		2nd toss					
		1	2	3	4	5	6
1st toss	1	0	1	2	3	4	5
	2	-1	0	1	2	3	4
	3	-2	-1	0	1	2	3
	4	-3	-2	-1	0	1	2
	5	-4	-3	-2	-1	0	1
	6	-5	-4	-3	-2	-1	0

$P[X=0] = \frac{6}{36}$
 $P[X=1] = \frac{5}{36} = P[X=-1]$
 $P[X=2] = \frac{4}{36} = P[X=-2]$
 $P[X=3] = \frac{3}{36} = P[X=-3]$
 $P[X=4] = \frac{2}{36} = P[X=-4]$
 $P[X=5] = \frac{1}{36} = P[X=-5]$

$$P[X=k] = \frac{6-|k|}{36}, \quad |k| \leq 5$$


3.3 Expected Value and Moments of Discrete Random Variable

3.21 $E[X] = 0 \cdot \frac{1}{16} + 1 \cdot \frac{8}{16} + 2 \cdot \frac{7}{16} = \frac{22}{16}$

$E[Y] = 0 \cdot \frac{1}{4} + 1 \cdot \frac{1}{2} + 2 \cdot \frac{1}{4} = 1$ which is much less than $E[X]$.

We will use $\text{VAR}[X] = E[X^2] - E[X]^2$:

$E[X^2] = 1 \cdot \frac{8}{16} + 4 \cdot \frac{7}{16} = \frac{36}{16}$

$E[Y^2] = 1 \cdot \frac{1}{2} + 4 \cdot \frac{1}{4} = \frac{3}{2}$

$\text{VAR}[X] = \frac{36}{16} - \left(\frac{22}{16}\right)^2 = \frac{82}{256}$

$\text{VAR}[Y] = \frac{3}{2} - 1^2 = \frac{1}{2}$

X has lower variance than Y.

3.22

(a) $E[X] = 1 \cdot \frac{12}{25} + 2 \cdot \frac{6}{25} + 3 \cdot \frac{4}{25} + 4 \cdot \frac{3}{25} = \frac{48}{25} = 1.92$
 $E[X^2] = 1 \cdot \frac{12}{25} + 4 \cdot \frac{6}{25} + 9 \cdot \frac{4}{25} + 16 \cdot \frac{3}{25} = \frac{120}{25}$
 $VAR[X] = \frac{120}{25} - \left(\frac{48}{25}\right)^2 = \frac{696}{625} = 1.114$

(b) $E[X] = 1 \cdot \frac{8}{15} + 2 \cdot \frac{4}{15} + 3 \cdot \frac{2}{15} + 4 \cdot \frac{1}{15} = \frac{26}{15} = 1.73$
 $E[X^2] = 1 \cdot \frac{8}{15} + 4 \cdot \frac{4}{15} + 9 \cdot \frac{2}{15} + 16 \cdot \frac{1}{15} = \frac{58}{15}$
 $VAR[X] = \frac{58}{15} - \left(\frac{26}{15}\right)^2 = \frac{194}{225} = 0.862$

(c) $E[X] = 1 \cdot \frac{64}{105} + 2 \cdot \frac{32}{105} + 3 \cdot \frac{8}{105} + 4 \cdot \frac{1}{105} = \frac{156}{105} = 1.48$
 $E[X^2] = 1 \cdot \frac{64}{105} + 4 \cdot \frac{32}{105} + 9 \cdot \frac{8}{105} + 16 \cdot \frac{1}{105} = \frac{280}{105}$
 $VAR[X] = \frac{280}{105} - \left(\frac{156}{105}\right)^2 = \frac{5064}{(105)^2} = 0.459$

The means and variances decrease as we progress through these distributions.

3.23

$$E[X] = \sum_{i=0}^{15} i \frac{1}{16} = \frac{1}{16} \sum_{i=0}^{15} i$$

Let $S = 1 + 2 + \dots + k$

$$S = k + (k-1) + \dots + 1$$

$$\frac{2S}{2} = k(k+1)$$

$$\Rightarrow S = \sum_{i=1}^k i = \frac{k(k+1)}{2}$$

$$\therefore E[X] = \frac{1}{16} \sum_{i=1}^{15} i = \frac{1}{16} \frac{15(16)}{2} = \frac{15}{2} = 7.5$$

$$E[X^2] = \frac{1}{16} \sum_{i=1}^{15} i^2 \quad \sum_{i=1}^k i^2 = \frac{k(k+1)(2k+1)}{6}$$

$$= \frac{1}{16} \frac{15(16)(31)}{6} = \frac{155}{2}$$

$$\text{VAR}[X] = \frac{155}{2} - \left(\frac{15}{2}\right)^2 = \frac{310 - 225}{4} = \frac{85}{4}$$

3.24

$$E[X] = 2P[X=2] + 3P[X=3] + 4P[X=4]$$

$$= 2 \cdot \frac{1}{2} + 3 \cdot \frac{1}{4} + 4 \cdot \frac{1}{4} = 2\frac{3}{4} \text{ bits/block}$$

Let X_1, X_2, \dots, X_n be the codeword lengths for a sequence of source outputs. The average codeword

length is

$$\frac{1}{n} \sum_{i=1}^n X_i \rightarrow E[X] \text{ for large } n$$

$\therefore E[X]$ is the long-term average number of bits per block.

3.25 Without replacement

$$E[X] = 2 \cdot \frac{4}{5} + 51 \cdot \frac{1}{5} = \frac{59}{5} = 11.80$$

$$E[X^2] = 4 \cdot \frac{4}{5} + 51^2 \cdot \frac{1}{5} = \frac{2617}{5}$$

$$\text{VAR}[X] = \frac{2617}{5} - \left(\frac{59}{5}\right)^2 = \frac{9604}{25} = 384.16$$

with replacement:

$$E[X] = 2 \cdot \frac{81}{100} + 51 \cdot \frac{18}{100} + 100 \cdot \frac{1}{100} = \frac{1180}{100} = 11.80$$

$$E[X^2] = 4 \cdot \frac{81}{100} + 51^2 \cdot \frac{18}{100} + 10^4 \cdot \frac{1}{100} = \frac{57142}{100}$$

$$\text{VAR}[X] = \frac{57142}{100} - \left(\frac{1180}{100}\right)^2 = \frac{43218}{100} = 432.18$$

Means in both draws is the same!

3.26

$$E[Y] = \sum_{j=-5}^5 j P[Y=j]$$
$$= -5 \cdot \frac{1}{36} + 4 \cdot \frac{2}{36} - 3 \cdot \frac{3}{36} - 2 \cdot \frac{4}{36} - 1 \cdot \frac{5}{36} + 0 \cdot \frac{6}{36}$$
$$+ 1 \cdot \frac{5}{36} + 2 \cdot \frac{4}{36} + 3 \cdot \frac{3}{36} + 4 \cdot \frac{2}{36} + 5 \cdot \frac{1}{36}$$
$$= 0$$
$$\text{VAR}[Y] = E[Y^2] = \sum_{j=-5}^5 j^2 P[Y=j]$$
$$= \sum_{j=1}^5 j^2 [P[X=j] + P[X=-j]]$$
$$= 1 \cdot \frac{10}{36} + 4 \cdot \frac{8}{36} + 9 \cdot \frac{6}{36} + 16 \cdot \frac{4}{36} + 25 \cdot \frac{2}{36}$$
$$= \frac{185}{36}$$

3.27 $E[X] = \sum_{j=1}^{\infty} j P[X=j] = \sum_{j=1}^{\infty} j \frac{c}{j^2} = c \sum_{j=1}^{\infty} \frac{1}{j} = \infty$

mean does not exist.

$E[X^2] = \sum_{j=1}^{\infty} j^2 \frac{c}{j^2} = c \sum_{j=1}^{\infty} 1 = \infty$

none of the moments exist

This pmf decays sufficiently fast that probabilities add to 1, but too slowly for moments to exist.

3.28 $E[Y] = -1 \cdot \frac{1}{10} + 0 \cdot \frac{2}{10} + 1 \cdot \frac{3}{10} + 2 \cdot \frac{4}{10} = \frac{10}{10} = 1$

$E[Y^2] = 1 \cdot \frac{1}{10} + 1 \cdot \frac{3}{10} + 4 \cdot \frac{4}{10} = \frac{20}{10} = 2$

$\text{VAR}[Y] = 2 - 1^2 = 1.$

3.29 $P[X=j] = \left(\frac{1}{2}\right)^j$

$$E[X] = \sum_{j=1}^{\infty} j \left(\frac{1}{2}\right)^j = \frac{1}{2} \sum_{j=1}^{\infty} j \left(\frac{1}{2}\right)^{j-1} = \frac{1}{2} \sum_{j=0}^{\infty} j \left(\frac{1}{2}\right)^j$$

From geometric series we have

$$\sum_{j=0}^{\infty} \alpha^j = \frac{1}{1-\alpha}$$

$$\therefore \frac{d}{d\alpha} \sum_{j=0}^{\infty} \alpha^j = \sum_{j=0}^{\infty} j \alpha^{j-1} = \frac{1}{(1-\alpha)^2}$$

$$\therefore \sum_{j=0}^{\infty} j \left(\frac{1}{2}\right)^{j-1} = \frac{1}{\left(1-\frac{1}{2}\right)^2} = 4$$

and $E[X] = \frac{1}{2} \cdot 4 = 2$

3.30 From problem 3.19 a 5-bit codeword is decoded erroneously with probability $P_e = 0.00856$.

In 1000 ^{codeword} transmissions we expect only 8.56 to be in error.

In 1000 single bit transmissions, since $p = \frac{1}{10}$ we expect $1000 \cdot \frac{1}{10} = 100$ to be in error.

\therefore Error rate is reduced at expense of slower information transmission rate.

3.31

$$P[X=k] = \binom{n}{k} \left(\frac{1}{2}\right)^n$$

$$E[aX^2 + bX] = aE[X^2] + bE[X]$$

$$E[X] = \sum_{j=0}^n j \binom{n}{j} \left(\frac{1}{2}\right)^n = \left(\frac{1}{2}\right)^n \sum_{j=0}^n j \frac{n!}{j!(n-j)!}$$

$$= \left(\frac{1}{2}\right)^n \sum_{j=1}^n \frac{n!}{(j-1)!(n-j)!} \quad \text{let } j' = j-1$$

$$= \left(\frac{1}{2}\right)^n n \sum_{j'=0}^{n-1} \frac{(n-1)!}{j'!(n-1-j')!} = n \left(\frac{1}{2}\right)^n \sum_{j'=0}^{n-1} \binom{n-1}{j'}$$

$$= n \left(\frac{1}{2}\right)^n 2^{n-1} = \frac{n}{2}$$

$$E[X^2] = \sum_{j=0}^n j^2 \binom{n}{j} \left(\frac{1}{2}\right)^n = \left(\frac{1}{2}\right)^n n \sum_{j=1}^n j \frac{(n-1)!}{(j-1)!(n-j)!}$$

$$= n \left(\frac{1}{2}\right)^n \sum_{j'=0}^{n-1} (j'+1) \binom{n-1}{j'}$$

$$= n \left(\frac{1}{2}\right)^n \left[\underbrace{\sum_{j'=0}^{n-1} j' \binom{n-1}{j'} \left(\frac{1}{2}\right)^{n-1}}_{\substack{(n-1) \frac{1}{2} \\ \text{expected value of} \\ \text{binomial}}} + \underbrace{\sum_{j'=0}^{n-1} \binom{n-1}{j'} \left(\frac{1}{2}\right)^{n-1}}_1 \right]$$

binomial probs

$$= \frac{n}{2} \left[\frac{n}{2} + 1 \right]$$

$$\therefore E[aX^2 + bX] = a \frac{n}{2} \left(\frac{n}{2} + 1 \right) + b \frac{n}{2} \quad \checkmark \quad \text{average reward.}$$

3.31b
$$E[a^X] = \sum_{j=0}^n a^j \binom{n}{j} \left(\frac{1}{2}\right)^j = \sum_{j=0}^n \binom{n}{j} \left(\frac{a}{2}\right)^j$$

$$= \left(1 + \frac{a}{2}\right)^n$$

3.32 (a) $E[I_A(X)] = E\left[\frac{I(X)}{A}\right] \quad A = \{X > 10\}$

$$= \sum_{i=1}^{15} \frac{I_A(i)}{A} P[X=i] = \sum_{i=11}^{15} P[X=i]$$

$$= p_1 \left[\sum_{i=11}^{15} \frac{1}{i} \right] = \frac{\sum_{i=11}^{15} 1/i}{\sum_{i=1}^{15} 1/i} = 0.1173$$

(b) $E\left[\frac{I_A(X)}{A}\right] = p_1 \frac{\sum_{i=11}^{15} 1/2^{(i-1)}}{\sum_{i=1}^{15} 1/2^{(i-1)}}$

$$= \frac{\sum_{i=11}^{15} 1/2^{(i-1)}}{\sum_{i=1}^{15} 1/2^{(i-1)}} = 0.00946$$

prob has faster decay than (a)

(c) $E\left[\frac{I_A(X)}{A}\right] = p_1 \frac{\sum_{i=11}^{15} \left(\frac{1}{2}\right)^{j(j-1)/2}}{\sum_{i=1}^{15} \left(\frac{1}{2}\right)^{j(j-1)/2}}$

$$= 1.69 \times 10^{-17}$$

The last prob decays much faster than the first two.

3.33

Ⓐ $E[(X-10)^+] = \sum_{i=1}^{15} (i-10) P[X=i] = p \sum_{i=1}^{15} (i-10) \frac{1}{2} = 0.33373$

Ⓑ $E[(X-10)^+] = p \sum_{i=1}^{15} (i-10) 2^{-(i-1)} = 0.00174$

Ⓒ $E[(X-10)^+] = p \sum_{i=1}^{15} (i-10) 2^{-i(i-1)/2} = 1.69 \times 10^{-17}$

3.34

$X_{\max} = m$ since casino has 2^m dollars.

Ⓐ \therefore casino can only play up to

Ⓑ $E[Y] = \sum_{k=1}^m 2^k \left(\frac{1}{2}\right)^k = m$

Ⓒ Player is willing to pay at most m dollars.

3.4 Conditional Probability Mass Function

3.35 (a) $P[X=k|X>0] = \begin{cases} 8/15 & k=1 \\ 7/15 & k=2 \end{cases}$ since $P[X=k|X>0] = \frac{P[X=k]}{P[X>0]}$

(b) $P[X=k|N_m=1] = \frac{P[X=k, N_m=1]}{P[N_m=1]}$

$= \begin{cases} \frac{\frac{1}{4} \cdot \frac{1}{2} + \frac{1}{2} \cdot \frac{1}{2}}{\frac{1}{2}} = \frac{3}{4} & k=1 \text{ Michael} \\ \frac{\frac{1}{4} \cdot \frac{1}{2}}{\frac{1}{2}} = \frac{1}{4} & k=2 \end{cases}$

			values	
	0	1	2	
0				1/4
1	1	1	2	1/2
2				1/4
	1/4	1/2	1/4	

3.35 (c) We need to change underlying sample space

				00	01	10	11	
M	00							1/4
	01							1/4
	10	1	1	1	2			1/4
	11	2	2	2	2			1/4
	1/4	1/4	1/4	1/4				

$P[X=1 | M \in \{10, 11\}] = \frac{3}{16} / \frac{1}{2} = \frac{3}{16}$

$P[X=2 | M \in \{10, 11\}] = \frac{5}{16} / \frac{1}{2} = \frac{5}{16}$

(d) $P[N_c=2 | X=2] = \frac{P[N_c=2, X=2]}{P[X=2]}$

			N_c		
	0	1	2		
N_m	0			2	1/4
	1			2	1/2
	2	2	2	2	1/4
	1/4	6/16	9/16		

$= \frac{9}{16} / \left(\frac{1}{4} \left(\frac{1}{16} \right) + \frac{1}{4} \frac{6}{16} + \frac{9}{16} \right) = \frac{36}{43}$

3.36

(a)
$$P[X=k|X<4] = \frac{P[X=k]}{1 - P[X=4]} = \begin{cases} \frac{12}{22} & k=1 \\ \frac{4}{22} & k=2 \\ \frac{4}{22} & k=3 \end{cases}$$

(b)
$$P[X=k|X<4] = \frac{P[X=k]}{104/105} = \begin{cases} \frac{64}{104} & k=1 \\ \frac{32}{104} & k=2 \\ \frac{8}{104} & k=3 \end{cases}$$

(c)
$$P[X=k|X<4] = \frac{P[X=k]}{14/15} = \begin{cases} \frac{8}{14} & k=1 \\ \frac{4}{14} & k=2 \\ \frac{2}{14} & k=3 \end{cases}$$

3.37

(a)
$$P[X=k|X<8] = \frac{P[X=k]}{P[X<8]} = \frac{1/16}{1/2} = \frac{1}{8} \text{ for } k < 8$$

(b)
$$P[X=k | \text{1st bit is zero}] = P[X=k | X < 8] \text{ same as (a).}$$

(c)
$$P[X=k | \text{4th bit is zero}] = P[X=k | X \text{ is even}]$$

$$= \frac{P[X=k, k \text{ even}]}{P[X \text{ is even}]} = \frac{1/16}{1/2} = \frac{1}{8} \text{ for } k \text{ even}$$

3.38) "No message gets through" $\Leftrightarrow X > 1$

(a)
$$P[X=k | X > 1] = \frac{P[X=k]}{P[X > 1]} = \frac{(\frac{1}{2})^k}{\frac{1}{2}} = (\frac{1}{2})^{k-1} \text{ for } k > 1$$

(b) If 1st transmitter transmitted in slot 1, then
 collision occurs w/ time slot 1 with prob $\frac{1}{2} \Leftrightarrow X > 1$
 success " " " " " " $\frac{1}{2} \Leftrightarrow X = 1$

$\therefore P[X=1 | C] = \frac{1}{2}$

for $k > 1$

$$P[X=k | C] = P[X=k, X > 1] = P[X=k | X > 1] P[X > 1]$$

$$= (\frac{1}{2})^{k-1} \cdot \frac{1}{2} = (\frac{1}{2})^k \quad k > 1$$

\therefore knowledge that C occurred does not change the pmf of X.

3.39

$$\begin{aligned}
 \textcircled{a} P[X=k | X > 1] &= \left(\frac{1}{2}\right)^{k-1} \quad k=2,3,\dots \\
 E[X | X > 1] &= \sum_{k=2}^{\infty} k \left(\frac{1}{2}\right)^{k-1} = \sum_{k'=1}^{\infty} (k'+1) \left(\frac{1}{2}\right)^{k'} \quad \text{where } k'=k-1 \\
 &= \sum_{k'=0}^{\infty} k' \left(\frac{1}{2}\right)^{k'} + \sum_{k'=1}^{\infty} \left(\frac{1}{2}\right)^{k'} \\
 &= \underbrace{E[X]}_{\text{avg. \# as of starting from scratch}} + \underbrace{1}_{\text{1 transmission is contain}} = 3
 \end{aligned}$$

3.39b

"message gets through w/ 1st time slot" = $X=1$

$$P[X=k|X=1] = \begin{cases} 0 & k > 1 \\ 1 & k=1 \end{cases}$$

$$E[X|X=1] = 1 \cdot P[X=1] = 1$$

(c) Let $A = \{X=1\}$ $B = \{X>1\}$ then $A \cup B$ form a partition

$$E[X] = E[X|A]P[A] + E[X|B]P[B]$$

$$= 1 \cdot \frac{1}{2} + 3 \cdot \frac{1}{2} = 2$$

note ~~we~~ we can use the result of part a to find $E[X]$:

$$E[X] = 1 \cdot \frac{1}{2} + (E[X]+1) \frac{1}{2} \Rightarrow E[X] = 2.$$

$$\begin{aligned} \text{(d)} \quad E[X^2|X>1] &= \sum_{k=2}^{\infty} k^2 \left(\frac{1}{2}\right)^{k-1} = \sum_{k'=1}^{\infty} (k'+1)^2 \left(\frac{1}{2}\right)^{k'} \\ &= \sum_{k'=1}^{\infty} k'^2 \left(\frac{1}{2}\right)^{k'} + 2 \sum_{k'=1}^{\infty} k' \left(\frac{1}{2}\right)^{k'} + \sum_{k'=1}^{\infty} \left(\frac{1}{2}\right)^{k'} \end{aligned}$$

$$= E[X^2] + 2E[X] + 1$$

$$= E[X^2] + 5$$

$$E[X^2|X=1] = 1$$

$$\begin{aligned} \therefore E[X^2] &= E[X^2|X=1] \frac{1}{2} + E[X^2|X>1] \frac{1}{2} \\ &= \frac{1}{2} + [E[X^2] + 5] \frac{1}{2} \end{aligned}$$

$$\Rightarrow E[X^2] = 6$$

$$\text{VAR}[X] = E[X^2] - E[X]^2 = 6 - 2^2 = 2$$

3.40

By (3.31b)

$$E[X^2] = \sum_{i=1}^n E[X^2|B_i] P[B_i] \quad \text{and} \quad E[X] = \sum_{i=1}^n E[X|B_i] P[B_i]$$

$$\text{VAR}[X] = E[X^2] - E[X]^2$$

$$= \sum_{i=1}^n E[X^2|B_i] P[B_i] - \left(\sum_{i=1}^n E[X|B_i] P[B_i] \right)^2$$

$$\neq \sum_{i=1}^n (E[X^2|B_i] - E[X|B_i]^2) P[B_i]$$

3.41

$$(a) \quad P[X=j | \text{1st draw} = k]$$

 $k = 1, 50$

$$P[X=j | \text{1st draw} = 1] = \begin{cases} \frac{8}{9} & j=2 \\ \frac{1}{9} & j=51 \end{cases}$$

$$P[X=j | \text{1st draw} = 50] = \begin{cases} 1 & j=51 \\ 0 & \text{otherwise} \end{cases}$$

$$(b) \quad E[X | \text{1st draw} = 1] = 2 \cdot \frac{8}{9} + 51 \cdot \frac{1}{9}$$

$$E[X | \text{1st draw} = 50] = 51$$

$$(c) \quad E[X] = E[X|1] \cdot \frac{9}{10} + E[X|50] \cdot \frac{1}{10}$$

$$= \frac{67}{9} \cdot \frac{9}{10} + \frac{51}{10} = \frac{118}{10}$$

$$(d) \quad E[X^2 | 1] = 4 \cdot \frac{8}{9} + (51)^2 \cdot \frac{1}{9} \quad E[X^2 | 50] = (51)^2$$

$$E[X^2] = \left(4 \cdot \frac{8}{9} + \frac{51^2}{9} \right) \frac{9}{10} + \frac{51^2}{10} = \frac{32}{10} + 2 \left(\frac{51^2}{10} \right) = \frac{5234}{10}$$

$$\text{VAR}[X] = \frac{32}{10} + 2 \left(\frac{51^2}{10} \right) - \left(\frac{118}{10} \right)^2 = 384.16$$

3.42 Assume # of heads is k

then $E[Y|k] = n - 2k$

$$\therefore E[Y] = \sum_{k=0}^n E[Y|k]P[k] = \sum_{k=0}^n (n - 2k) \binom{n}{k} p^k (1-p)^{n-k}$$

$$= n - 2E[X] = n - 2np$$

$$= n(1 - 2p)$$

Similarly

$$E[Y^2|k] = (n - 2k)^2 = n^2 - 4kn + 4k^2$$

$$E[Y^2] = \sum_{k=0}^n (n^2 - 4kn + 4k^2) \binom{n}{k} p^k (1-p)^{n-k}$$

$$= n^2 - 4nE[X] + 4E[X^2]$$

$$= n^2 - 4n^2p + 4(npq + (np)^2)$$

$$= n^2 - 4n^2p + 4npq + 4n^2p^2$$

$$\text{VAR}[Y] = E[Y^2] - E[Y]^2$$

$$= n^2 - 4n^2p + 4npq + 4n^2p^2 - \underbrace{n^2(1 - 2p)^2}_{1 - 4p + 4p^2}$$

$$= 4npq$$

3.43

ⓐ) If password has not been found after k tries then there remain $2^m - k$ possible passwords.

$$P[X=j | X > k] = \begin{cases} \frac{1}{2^m - k} & j = k+1, \dots, 2^m \\ 0 & \text{otherwise} \end{cases}$$

$$\begin{aligned} \text{ⓑ) } E[X | X > k] &= \sum_{j=k+1}^{2^m} j \frac{1}{2^m - k} = \frac{1}{2^m - k} \sum_{j=k+1}^{2^m} j \\ &= \frac{1}{2^m - k} \left[\frac{2^m(2^m + 1)}{2} - \frac{k(k+1)}{2} \right] \\ &= \frac{1}{2^m - k} \left[\frac{(2^m - k)(2^m + k + 1)}{2} \right] = \frac{2^m + k + 1}{2} \\ &= \underbrace{(k+1)}_{\text{minimum}} + \underbrace{\frac{2^m - (k+1)}{2}}_{\text{average additional number of tries}} \end{aligned}$$

3.5 Important Discrete Random Variables

3.44

(a) $S = \{1, 3, 3, 4, 5\}$ $A = \{\xi > 3\}$

$P[I_A = 0] = \frac{3}{5}$ $P[I_A = 1] = \frac{2}{5}$

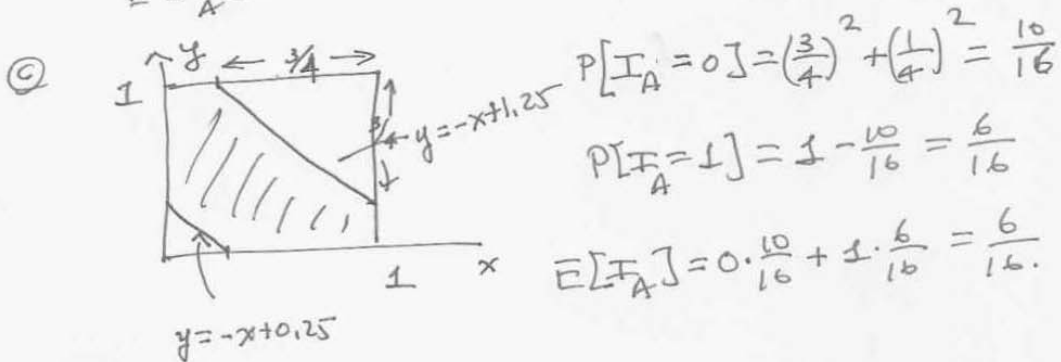
$E[I_A] = 0 \cdot \frac{3}{5} + 1 \cdot \frac{2}{5} = \frac{2}{5}$

(b) $S = [0, 1]$ $A = \{0.3 < \xi \leq 0.7\}$

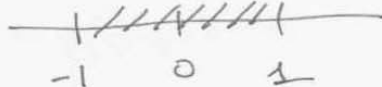
$P[I_A = 0] = P[\xi \leq 0.3] + P[0.7 < \xi \leq 1] = 0.6$

$P[I_A = 1] = P[0.3 < \xi \leq 0.7] = 0.4$

$E[I_A] = 0 \cdot 0.6 + 1 \cdot 0.4 = 0.4$



(d) $A = \{\xi > a\}$



$a < -1$

$P[I_A = 1] = 1$
 $E[I_A] = 1$

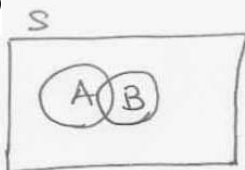
$-1 < a < 1$

$P[I_A = 1] = \frac{1-a}{2}$
 $E[I_A] = \frac{1-a}{2}$

$1 < a$

$P[I_A = 1] = 0$
 $E[I_A] = 0$

3.45



(a) $I_S = 1 \iff \xi \in S \Rightarrow I_S = 1 \text{ all } \xi$
 $I_\phi = 1 \iff \xi \in \phi \Rightarrow I_\phi = 0 \text{ all } \xi$

(b) $I_{A \cap B}(\xi) = 1 \iff \xi \in A \text{ and } \xi \in B \Leftrightarrow I_A(\xi) = 1 \text{ and } I_B(\xi) = 1$
 $\Leftrightarrow I_{A \cap B}(\xi) = I_A(\xi) I_B(\xi).$

$I_{A \cup B}(\xi) = 0 \iff \xi \notin A \cup B \Leftrightarrow \xi \in A^c \cap B^c \Leftrightarrow I_{A^c}(\xi) I_{B^c}(\xi) = 1$
 $\Leftrightarrow (1 - I_A(\xi))(1 - I_B(\xi)) = 1$
 $\Leftrightarrow 1 - I_A(\xi) - I_B(\xi) + I_A(\xi) I_B(\xi) = 1$
 $\Leftrightarrow I_A(\xi) + I_B(\xi) - I_A(\xi) I_B(\xi) = 0$
 $\Leftrightarrow I_A(\xi) + I_B(\xi) - I_{A \cap B}(\xi) = 0 = I_{A \cup B}(\xi)$

(c) $E[I_S] = 1 \cdot P[S] = 1$
 $E[I_\phi] = 0 \cdot P[\phi] = 0$
 $E[I_{A \cap B}] = 1 \cdot P[A \cap B]$
 $E[I_{A \cup B}] = E[I_A] + E[I_B] - E[I_{A \cap B}]$
 $= P[A] + P[B] - P[A \cap B].$

3.46 $n=8$ $p=0.25$ Binomial random variable

(a) $P[N=0] = \binom{8}{0} p^0 (1-p)^{8-0} = (0.75)^8 = 0.100$

(b) $P[N=1] = \binom{8}{1} p^1 (1-p)^{8-1} = 8(0.25)(0.75)^7 = 0.267$

(c) $P[N > 4] = \sum_{j=5}^8 \binom{8}{j} (0.25)^j (0.75)^{8-j} = 0.0273$

(d) $P[2 < N < 6] = \sum_{j=3}^5 \binom{8}{j} (0.25)^j (0.75)^{8-j} = 0.3172$

3.47

(a) $A_i = \{U_i < 0.25\}$
 $A_i^c = \{U_i > 0.25\}$

$P[A_1 A_2 A_3 A_4 A_5^c A_6^c A_7^c A_8^c] = (0.25)^4 (0.75)^4$
 = 0.00124

(b) $P[N=4] = \binom{8}{4} (0.25)^4 (0.75)^4 = 0.0865$

(c) $A_i = \{U_i < 0.25\}$
 $B_i = \{0.25 < U_i < 0.75\}$
 $C_i = \{U_i > 0.75\}$
 $P[A_1 A_2 A_3 B_4 B_5 C_6 C_7 C_8] = (0.25)^3 (0.5)^2 (0.25)^3$
 $= (0.25)^6 (0.5)^2$
 $= 6.10 \times 10^{-5}$

(d) $P[N_1=3, N_2=2, N_3=3] = \frac{8!}{3! 2! 3!} (0.25)^3 (0.5)^2 (0.25)^3$
 multinomial

(e) $P[A_1 A_2 A_3 A_4 C_5 C_6 C_7 C_8] = (0.25)^4 (0.25)^4 = 1.526 \times 10^{-5}$

(f) $P[N_1=4, N_2=0, N_3=4] = \frac{8!}{4! 0! 4!} (0.25)^4 (0.5)^0 (0.25)^4$
 $= 0.00107$

3.48

This Octave program will plot binomial pmf.
 $> n=4;$
 $> x=[0:n];$
 $> p=0.10;$
 $> stem(\text{binomial_pmf}(x, n, p))$

3.49

3.32 a) Let I_k denote the outcome of the k th Benoulli trials. The probability that the single event occurred in the k th trial is:

$$\begin{aligned} P\{I_k = 1|X = 1\} &= \frac{P\{I_k = 1 \text{ and } I_j = 0 \text{ for all } j \neq k\}}{P\{X = 1\}} \\ &= \frac{P[0 \ 0 \dots 1 \ 0 \dots 0]}{P\{X = 1\}} \\ &= \frac{p(1-p)^{n-1}}{\binom{n}{1} p(1-p)^{n-1}} = \frac{1}{n} \end{aligned}$$

Thus the single event is equally likely to have occurred in any of the n trials.

b) The probability that the two successes occurred in trials j and k is:

$$P\{I_j = 1, I_k = 1|X = 2\} = \frac{P\{I_j = 1, I_k = 1, I_m = 0 \text{ for all } m \neq j, k\}}{P\{X = 2\}}$$

3.50 a)
$$\frac{p_k}{p_{k-1}} = \frac{\binom{n}{k} p^k q^{n-k}}{\binom{n}{k-1} p^{k-1} q^{n-k+1}} = \frac{\frac{n!}{k!(n-k)!} p}{\frac{n!}{(k-1)!} q} = \frac{(n-k+1)p}{kq}$$

$$= \frac{(n+1)p - k(1-q)}{kq} = 1 + \frac{(n+1)p - k}{kq}$$

b) First suppose $(n+1)p$ is not an integer, then
 for $0 \leq k \leq [(n+1)p] < (n+1)p$

$$(n+1)p - k > 0$$

so

$$\frac{p_k}{p_{k-1}} = 1 + \frac{(n+1)p - k}{kq} > 1$$

$\Rightarrow p_k$ increases as k increases from 0 to $[(n+1)p]$
 for $k > (n+1)p \geq [(n+1)p]$

$$(n+1)p - k < 0$$

so

$$\frac{p_k}{p_{k-1}} = 1 + \frac{(n+1)p - k}{kq} < 1$$

$\Rightarrow p_k$ decreases as k increases beyond $[(n+1)p]$
 $\therefore p_k$ attains its maximum at $k_{MAX} = [(n+1)p]$
 If $(n+1)p = k_{MAX}$ then above implies that

$$\frac{p_{k_{MAX}}}{p_{k_{MAX}-1}} = 1 \Rightarrow p_{k_{MAX}} = p_{k_{MAX}-1}$$

3.51

$$\begin{aligned}
 \text{(a)} \quad (a+b+c)^n &= \sum_{k=0}^n \binom{n}{k} (a+b)^k c^{n-k} \\
 &= \sum_{k=0}^n \binom{n}{k} c^{n-k} \sum_{j=0}^k \binom{k}{j} a^j b^{k-j} \\
 &= \sum_{k=0}^n \sum_{j=0}^k \frac{n!}{k!(n-k)! j!(k-j)!} a^j b^{k-j} c^{n-k} \\
 &= \sum_{k=0}^n \sum_{j=0}^k \frac{n!}{j!(n-k)!(k-j)!} a^j b^{k-j} c^{n-k} \\
 &= \sum_{j_1, j_2, j_3} \frac{n!}{j_1! j_2! j_3!} a^{j_1} b^{j_2} c^{j_3} \\
 &\quad j_i \geq 0 \\
 &\quad j_1 + j_2 + j_3 = n
 \end{aligned}$$

$$\text{(c)} \quad 1 = (p_1 + p_2 + p_3)^n = \sum_{j_1, j_2, j_3} \frac{n!}{j_1! j_2! j_3!} p_1^{j_1} p_2^{j_2} p_3^{j_3}$$

3.52

$p = 0.01$ $N = \# \text{ errors detected until first error}$

$$\text{(a)} \quad P[N=k] = (1-p)^k p \quad k=0, 1, 2, \dots$$

$$\text{(b)} \quad E[N] = \sum_{k=0}^{\infty} k (1-p)^k p = (1-p)p \sum_{k=0}^{\infty} k (1-p)^{k-1}$$

$$= (1-p)p \frac{1}{(1-(1-p))^2} = \frac{1-p}{p} \quad \text{by Eqn 3.14}$$

$$\text{(c)} \quad 0.99 = P[N > k_0] = \sum_{k=k_0+1}^{\infty} (1-p)^k p = p (1-p)^{k_0+1} \sum_{k=0}^{\infty} (1-p)^k$$

$$= (1-p)^{1001} \Rightarrow p = 1 - 0.99^{\frac{1}{1001}} = 4.004 \times 10^{-5}$$

3.53 N geometric $n = 1, 2, \dots$

(a)
$$P[N=k | N \leq m] = \frac{P[N=k, N \leq m]}{P[N \leq m]} = \frac{P[N=k]}{P[N \leq m]} \quad 1 \leq k \leq m$$

$$= \frac{p(1-p)^{k-1}}{\sum_{j=1}^m p(1-p)^{j-1}} = \frac{p(1-p)^{k-1}}{1-(1-p)^m} \quad 1 \leq k \leq m$$

(b)
$$P[N \text{ odd}] = \sum_{j=0}^{\infty} p(1-p)^{2j+1} = p(1-p) \sum_{j=0}^{\infty} ((1-p)^2)^j$$

$$= \frac{p(1-p)}{1-(1-p)^2}$$

3.54
$$P[M \geq k+j | M > j] = \frac{P[M \geq k+j, M > j]}{P[M > j]} = \frac{P[M \geq k+j]}{P[M > j]} \quad \text{for } k \geq 1$$

$$= \frac{\sum_{i=k+j}^{\infty} p(1-p)^{i-1}}{\sum_{i=j+1}^{\infty} p(1-p)^{i-1}}$$

$$= \frac{(1-p)^{k+j-1}}{(1-p)^j} = (1-p)^{k-1} = P[M \geq k]$$

The probability of k additional trials until the first success is independent of how many failures have already transpired.

3.55

3.36 The memoryless property states that for $j, k \geq 1$.

$$\begin{aligned} P[M \geq k] &= P[M \geq k + j | M > j] \\ &= \frac{P[M \geq k + j]}{P[M > j]} = \frac{P[M \geq k + j]}{P[M \geq j + 1]} \end{aligned}$$

\Rightarrow

$$P[M \geq k + j] = P[M \geq k]P[M \geq j + 1]$$

Let

$$a_k = P[M \geq k],$$

then we have

$$(*) \quad a_{k+j} = a_k a_{j+1} \quad j \geq 1, k \geq 1$$

where $a_1 = 1$ and $a_2 = 1 - P[M = 1] = 1 - p$.

Equation (*) with $j = 1$ becomes

$$a_{k+1} = a_2 a_k \quad k \geq 1$$

$$\Rightarrow a_k = a_2^{k-1} \quad k \geq 1$$

$$\Rightarrow P[M \geq k] = (1 - p)^{k-1} \quad k \geq 1$$

$$\begin{aligned} P[M = k] &= P[M \geq k] - P[M \geq k + 1] \\ &= (1 - p)^{k-1} - (1 - p)^k \\ &= (1 - p)^{k-1}(1 - (1 - p)) \\ &= (1 - p)^{k-1}p \end{aligned}$$

3.56 $C_{\text{rent}} = \$50$ $C_{\text{repair}} = \$20$ $p = \frac{1}{12}$ $N = \text{# of ads} \rightarrow 12 \text{ months}$

$$P[N=k] = \binom{12}{k} \left(\frac{1}{12}\right)^k \left(\frac{11}{12}\right)^{12-k} \quad k=0, 1, \dots, 12$$

$0.99 = \sum_{j=k_0}^{12} P[N=j] \Rightarrow k_0 = ?$

k	$P[N=k]$	$P[N \leq k]$
0	0.352	0.352
1	0.384	0.736
2	0.192	0.928
3	0.058	0.986

$\Rightarrow k_0 = 3$

$\Rightarrow \text{Change } \$50 + 3 \times \$20 = \$110.$

Avg cost per player = $\$50 + \$20 E[N] = \$70.$
 $E[N] = 12 \left(\frac{1}{12}\right) = 1$

3.57 $\alpha_s = 48$ $\alpha_r = 24$ $\alpha_g = 12$ $\text{slit} = \frac{1}{12}$

Ⓐ $P[N_g = 0] = \frac{(\alpha_g \frac{1}{12})^0}{0!} e^{-\alpha_g \frac{1}{12}} = e^{-1} = 0.368$

Ⓑ $P[N_g = 0, N_r \leq 2] = P[N_g = 0] P[N_r \leq 2] = e^{-1} \sum_{k=0}^2 \frac{(\alpha_r \frac{1}{12})^k}{k!} e^{-\alpha_r \frac{1}{12}}$

$$= e^{-1} \left[e^{-2} + \frac{2}{1!} e^{-2} + \frac{4}{2!} e^{-2} \right]$$

$$= e^{-3} [1 + 2 + 2] = 5e^{-3} = 0.249$$

Ⓒ $P[N_g = 0, N_r = 0, N_s \geq 5] = P[N_g = 0] P[N_r = 0] P[N_s \geq 5]$

$$= e^{-1} e^{-2} e^{-4} \sum_{k=0}^5 \frac{4^k}{k!} = e^{-3} (0.785)$$

$$= 0.039$$

It's hard to avoid the red and green thingsies!

3.58

$$P[X > 4] < 0.9 \Leftrightarrow P[X \leq 4] > 0.1$$

$$P[X \leq 4] = \sum_{k=0}^4 \frac{\alpha^k}{k!} e^{-\alpha} = \sum_{k=0}^4 \frac{(5/n)^k}{k!} e^{-5/n}$$

Since $\alpha = \frac{\lambda}{n\mu} = \frac{5}{n}$
 If $n = 2$ then $P[X \leq 4] = 0.811$. Therefore ^{two} ~~one~~ employees ^{are} ~~is~~ sufficient.

$$P[X = 0] = e^{-\alpha} = e^{-2.5} = 0.082$$

3.59

$\lambda = 6000 \text{ requests/minute} = 100 \text{ requests/sec}$
 $\alpha = \lambda \frac{1}{10} = 10 \text{ requests/100ms}$

a) $P[N=0] = e^{-10} = 4.54 \times 10^{-5}$

b) $P[5 \leq N \leq 10] = \sum_{k=5}^{10} \frac{10^k}{k!} e^{-10} = 0.554$

3.60 Use octave to plot pmf.

> a = 0.1 ;
 > j = [0:20];
 > stem (poisson_pdf (j, a))

3.61

$$\mathcal{E}[X] = \sum_{k=0}^{\infty} k \frac{\alpha^k}{k!} e^{-\alpha} = \alpha \sum_{k=1}^{\infty} \frac{\alpha^{k-1}}{(k-1)!} e^{-\alpha} = \alpha \underbrace{\sum_{k'=0}^{\infty} \frac{\alpha^{k'}}{k'!} e^{-\alpha}}_1 = \alpha$$

$$\mathcal{E}[X^2] = \sum_{k=0}^{\infty} k^2 \frac{\alpha^k}{k!} e^{-\alpha} = \alpha \sum_{k=1}^{\infty} k \frac{\alpha^{k-1}}{(k-1)!} e^{-\alpha}$$

$$= \alpha \sum_{k'=0}^{\infty} (k'+1) \frac{\alpha^{k'}}{k'!} e^{-\alpha} = \alpha \{ \alpha + 1 \}$$

$$\sigma_X^2 = \mathcal{E}[X^2] - \mathcal{E}[X]^2 = \alpha \{ \alpha + 1 - \alpha \}$$

$$= \alpha$$

3.62 $\frac{p_k}{p_{k-1}} = \frac{\frac{\alpha^k}{k!} e^{-\alpha}}{\frac{\alpha^{k-1}}{(k-1)!} e^{-\alpha}} = \frac{\alpha}{k}$

If $\alpha < 1$ then $\frac{p_k}{p_{k-1}} = \frac{\alpha}{k} < 1$ for $k \geq 1$
 $\therefore p_k$ decreases as k increases from 0
 $\therefore p_k$ attains its maximum at $k = 0$

If $\alpha > 1$ then
 for $0 \leq k \leq [\alpha] < \alpha$, $\frac{p_k}{p_{k-1}} = \frac{\alpha}{k} > 1$
 $\Rightarrow p_k$ increase from $k = 0$ to $k = [\alpha]$
 for $[\alpha] < \alpha < k$, $\frac{p_k}{p_{k-1}} = \frac{\alpha}{k} < 1$
 $\Rightarrow p_k$ decreases as k increases beyond $[\alpha]$
 $\therefore p_k$ attains its maximum at $k_{\max} = [\alpha]$

If $\alpha = [\alpha]$ then for $k = [\alpha]$

$$\frac{p_k}{p_{k-1}} = 1 \Rightarrow p_{k_{\max}} = p_{k_{\max} - 1}$$

3.63

$n = 10$	$p = 0.1$	$np = 1$		
	$k = 0$	$k = 1$	$k = 2$	$k = 3$
Binomial	0.3487	0.387	0.1937	0.0574
Poisson	0.3679	0.3679	0.1839	0.0613

$n = 20$	$p = 0.05$	$np = 1$		
	$k = 0$	$k = 1$	$k = 2$	$k = 3$
Binomial	0.3585	0.3774	0.1887	0.06
Poisson	0.3679	0.3679	0.1839	0.0613

$n = 100$	$p = 0.01$	$np = 1$		
	$k = 0$	$k = 1$	$k = 2$	$k = 3$
Binomial	0.366	0.3697	0.1849	0.061
Poisson	0.3679	0.3679	0.1839	0.0613

3.64 N Poisson $\lambda = \frac{3}{\cancel{100}}$ $R = 2 \times 10^6$ bps

@ $X = R/N$ "infinite" for $N=0$ R/k for $N=k \geq 1$

$S_x = \left\{ \infty, 20, 10, \frac{2}{3}, 0.5, \frac{1}{3}, \frac{2}{7}, \dots \right\}$

$P[X = R/k] = P[N = k]$

@ $0.9 = P[N \leq k] = \sum_{j=0}^k \frac{(3)^j}{j!} e^{-3}$ $P[N \leq 5] = 0.916$

↙ use octave
`poisson_cdf(k, 3)`

@ $X \geq 1 \Leftrightarrow k \leq 2$

$P[N \leq 2] = 0.423$

3.65 $n = 1000 \times 750 = 7.5 \times 10^5$ pixels

$p = 10^{-5}$ $np = 7.5$

$P[\text{display accepted}] = \sum_{k=0}^{15} \binom{n}{k} p^k (1-p)^{n-k} \approx \sum_{k=0}^{15} \frac{(7.5)^k e^{-7.5}}{k!}$

$= 0.9953$

3.66 $n = 10^4$ drives $p = 10^{-3}$ $np = 10^4(10^{-3}) = 10/\text{day}$

(a) $P[N=0] \approx e^{-np} = e^{-10} = 4.54 \times 10^{-5}$

(b) Failure rate in 2 days = 20
 $P[N \leq 10] = \sum_{j=0}^{10} \frac{(20)^j}{j!} e^{-20} = 1.08 \times 10^{-2}$

(c) $0.99 = P[N \leq k] = \sum_{j=0}^k \frac{10^j}{j!} e^{-10} \Rightarrow P[N \leq 17] = 0.986$

3.67 $p = 10^{-6}$ $n = 10^4$ $np = 10^{-2}$

(a) $P[N=0] = e^{-np} = 0.990$

$P[N \leq 3] = \sum_{k=0}^3 \frac{(0.01)^k}{k!} e^{-np} \approx 1$

(b) Find p so that
 $0.99 = P[N \geq 1] = 1 - P[N=0] = 1 - e^{-np}$
 $0.01 = e^{-np}$
 $\Rightarrow p = \frac{\ln 0.01}{n} = 4.6 \times 10^{-6}$

3.68

65 $E[X] = \sum_{k=1}^n kP[X=k] = \sum_{k=1}^n \frac{k}{n} = \frac{1}{n} \sum_{k=1}^n k = \frac{n(n+1)}{2n} = \frac{n+1}{2}$

$\sigma_X^2 = E[X^2] - E[X]^2 = \sum_{k=1}^n \frac{k^2}{n} - \left(\frac{n+1}{2}\right)^2 = \frac{(n+1)(2n+1)}{6} - \frac{(n+1)^2}{4}$

$= \frac{n^2 - 1}{12}$

3.69

X uniform in $\{-3, -2, \dots, 3, 4\}$ $P[X=j] = \frac{1}{8}$

(a) $E[X] = -4 + \frac{3+1}{2} = 0.5$

$VAR(X) = \frac{8^2-1}{12} = \frac{63}{12} = \frac{21}{4}$

(b) $E[Y] = E[-2X^2+3] = -2E[X^2]+3$
 $= -2[VAR(X)+E[X]^2]+3$

$= -2\left[\frac{21}{4} + (0.5)^2\right] + 3 = -8$

$E[Y^2] = E[(-2X^2+3)^2] = E[4X^4 - 12X^2 + 9]$

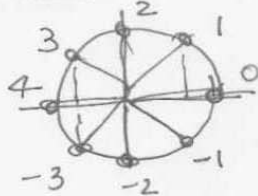
$VAR[Y] = E[Y^2] - E[Y]^2$

$= 4E[X^4] - 12E[X^2] + 9 - (-8)^2$

$E[X^4] = \frac{1}{8} [(-3)^2 + (-2)^2 + (-1)^2 + 0^2 + 1^2 + 2^2 + 3^2 + 4^2]$
 $= \frac{44}{8} = \frac{11}{2}$

$VAR[Y] = 4\left(\frac{11}{2}\right) - 12(-2) + 9 - 64 = 99$

(c) $W = \cos\left(\frac{\pi X}{8}\right)$



X	-3	-2	-1	0	1	2	3	4
W	$-\frac{1}{\sqrt{2}}$	0	$\frac{1}{\sqrt{2}}$	1	$\frac{1}{\sqrt{2}}$	0	$-\frac{1}{\sqrt{2}}$	-1

$P[W = -\frac{1}{\sqrt{2}}] = \frac{2}{8}$ $P[W = 0] = \frac{2}{8}$

$P[W = \frac{1}{\sqrt{2}}] = \frac{2}{8}$ $P[W = 1] = \frac{1}{8}$ $P[W = -1] = \frac{1}{8}$

3.70
$$p_k = \frac{1}{c_{10}} \frac{1}{k} \quad k=1, \dots, 10 \quad c_{10} = 2.93$$

$$p_1 = \frac{1}{2.93} = 0.3414$$

$$P[X > 5] = \frac{1}{c_{10}} \left[\frac{1}{6} + \dots + \frac{1}{10} \right] = 0.2204$$

3.71
$$p_k = \frac{1}{c_{1000}} \frac{1}{k} \quad c_{1000} = \ln 1000 + 0.57721 = 7.485$$

$$P[X \leq 10] = \frac{1}{c_{1000}} \sum_{j=1}^{10} \frac{1}{j} = \frac{c_{90}}{c_{1000}} = \frac{2.93}{7.485} = 0.3913$$

$$P[X > 990] = 1 - P[X \leq 990] = 1 - \frac{c_{990}}{c_{1000}}$$

$$= 1 - \frac{\ln 990 + 0.57721}{\ln 1000 + 0.57721} = 0.00134$$

3.72
$$p_k = \frac{1}{c_L} \frac{1}{k} \quad \ln p_k = \ln \frac{1}{c_L} + \ln \frac{1}{k}$$

$$= -\ln k - \ln c_L$$

$\ln p_k$ is linear in $\ln k$

3.73 $E[X] = \frac{L}{c_2} \approx \frac{L}{\ln L + 0.57721}$ for large L
 $VAR[X] = L^2/c_2^2 = E[X]^2$
 To plot $E[X]$ vs L use octave

```

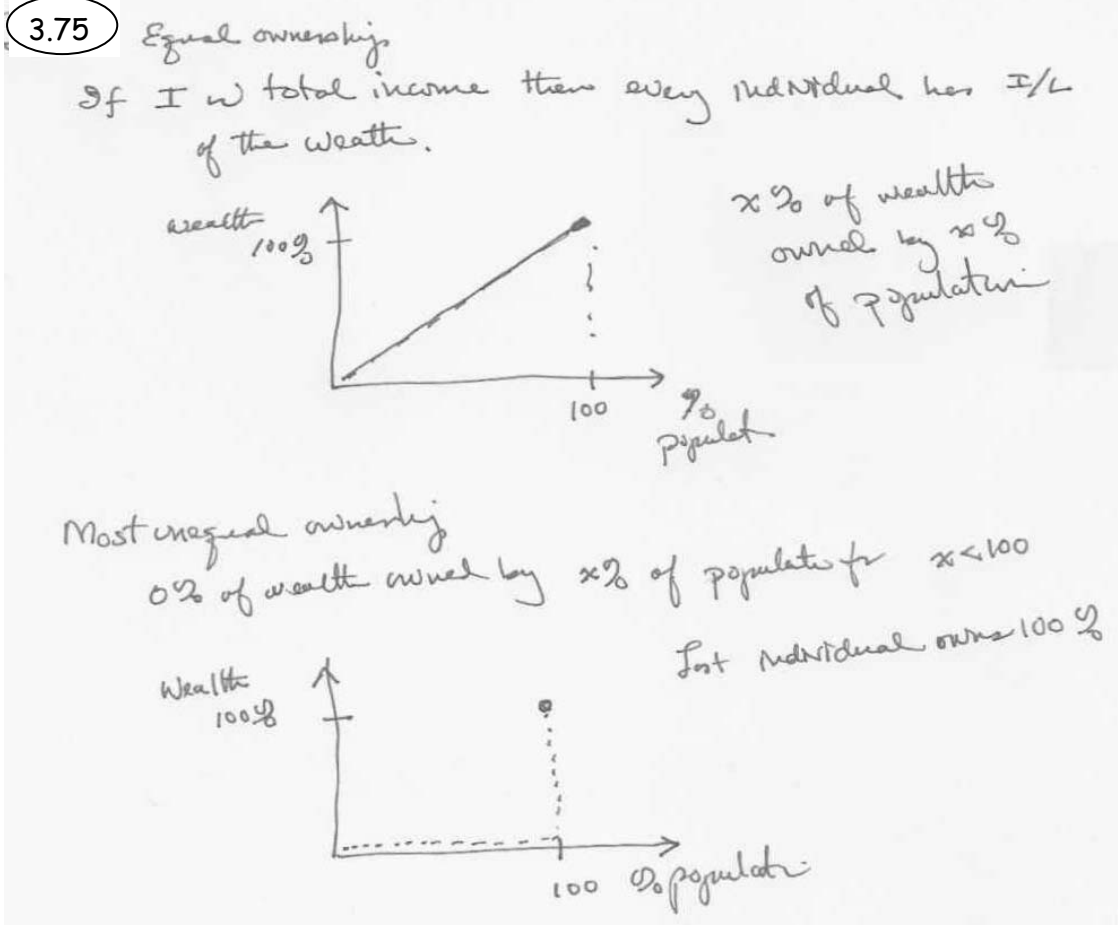
> L = [1:100];
> p = L.^(-1);
> cL = cumsum(p);          array of coefficients
> plot(L./cL)              plots means
> plot((L.^2)./(cL.^2))   plots variances
    
```

3.74 $p_k = \frac{1}{c_2} \frac{1}{k^2}$ $L = 10^4$ $c_2 = \ln 10^4 + 0.57721 = 9.7876$

$$0.99 = P[X \leq k_0] = \frac{1}{c_{10000}} \sum_{j=1}^{k_0} \frac{1}{j^2} = \frac{c_{k_0}}{c_{10000}} \approx \frac{\ln k_0 + 0.57721}{9.7876}$$

$$\ln k_0 \approx 0.99(9.7876) - 0.57721 = 9.067$$

Zipf decays very slowly!

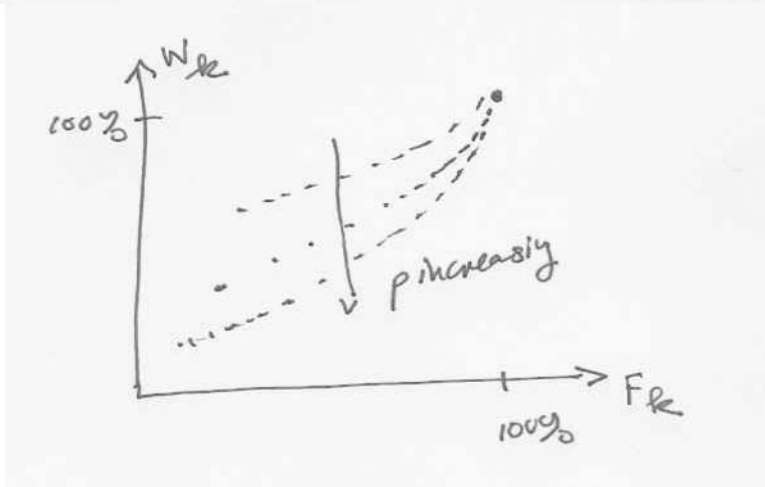


3.76 $P[X=k] = (1-p)p^{k-1} \quad k=1,2,3,\dots$

$F_k = P[X \leq k] = \sum_{j=1}^k p^{j-1} = 1-p^k$

$W_k = \frac{\sum_{j=1}^k j c p^{j-1}}{\sum_{j=1}^{\infty} j c p^{j-1}} = \frac{\sum_{j=1}^k j p^{j-1}}{\sum_{j=1}^{\infty} j p^{j-1}} = \frac{(1-p)^{-(k+1)} - (k+1)p(1-p)}{1-p^2}$

$\sum_{j=0}^m j p^j = \frac{d}{dp} \frac{1-p^{m+1}}{1-p} = \frac{(1-p^{m+1}) - (m+1)p^m(1-p)}{1-p^2}$



(3.77) $X = \frac{1}{z_\alpha} \frac{1}{k^\alpha} \quad k=1,2,\dots$

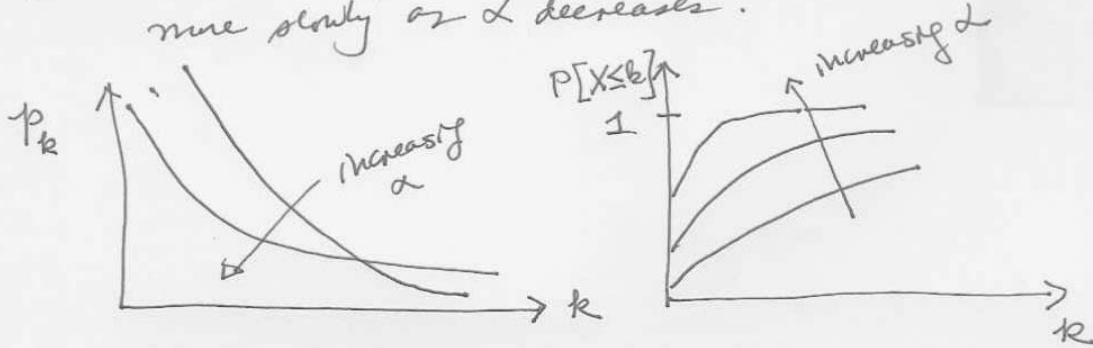
(a) $P[X \leq k] = \frac{1}{z_\alpha} \sum_{j=1}^k \frac{1}{j^\alpha} = \frac{z_{\alpha,k}}{z_\alpha}$ where $z_{\alpha,k} = \sum_{j=1}^k \frac{1}{j^\alpha}$

(b) z_α is given by the zeta function evaluated at α
 Using octave, we have
 > zeta(1.5)
 > ans = 2.6124

$z(1.5) = 2.6124$
 $z(2) = 1.6449$
 $z(3) = 1.2021$

If we add the first 100 terms to estimate z_α we have:
 $z_{1.5,100} = 2.41 \quad z_{2,100} = 1.635 \quad z_{3,100} = 1.202$

The series that defines the zeta function decays more slowly as α decreases.



3.6 Generation of Discrete Random Variables

3.78 The following Octave commands will give the requested plots:

(a)

```
x = [0:1:10];  
lambda = 0.5;  
figure;  
plot(x, poisson_pdf(x, lambda));  
figure;  
plot(x, poisson_cdf(x, lambda));  
figure;  
plot(x, 1-poisson_cdf(x, lambda));
```

```
x = [0:1:20];  
lambda = 5;  
figure;  
plot(x, poisson_pdf(x, lambda));  
figure;  
plot(x, poisson_cdf(x, lambda));  
figure;  
plot(x, 1-poisson_cdf(x, lambda));
```

```
x = [0:1:100];  
lambda = 50;  
figure;  
plot(x, poisson_pdf(x, lambda));  
figure;  
plot(x, poisson_cdf(x, lambda));  
figure;  
plot(x, 1-poisson_cdf(x, lambda));
```

(b)

```
x = [0:1:15];  
figure;  
plot(x, binomial_pdf(x, 48, 0.1));  
figure;  
plot(x, binomial_cdf(x, 48, 0.1));  
figure;  
plot(x, 1-binomial_cdf(x, 48, 0.1));
```

```
x = [0:1:30];  
figure;  
plot(x, binomial_pdf(x, 48, 0.3));  
figure;  
plot(x, binomial_cdf(x, 48, 0.3));  
figure;  
plot(x, 1-binomial_cdf(x, 48, 0.3));
```

```
x = [0:1:50];  
figure;
```

```
plot(x, binomial_pdf(x, 48, 0.5));  
figure;  
plot(x, binomial_cdf(x, 48, 0.5));  
figure;  
plot(x, 1-binomial_cdf(x, 48, 0.5));  
  
x = [20:1:50];  
figure;  
plot(x, binomial_pdf(x, 48, 0.75));  
figure;  
plot(x, binomial_cdf(x, 48, 0.75));  
figure;  
plot(x, 1-binomial_cdf(x, 48, 0.75));
```

(c)

```
x = [0:1:10];  
n = 100; p = 0.01;  
figure;  
plot(x, binomial_pdf(x, n, p), "1");  
hold on;  
plot(x, poisson_pdf(x, n*p), "3");  
hold off;
```

3.79

The following Octave commands produce the request plots:

(a)

```
L = 10;  
k = [1:1:L];  
cL = sum(1./k);  
pk = (1/cL)./k;  
figure;  
plot(k, discrete_pdf(k, k, pk));  
figure;  
plot(k, discrete_cdf(k, k, pk));  
figure;  
plot(k, 1-discrete_cdf(k, k, pk));
```

```
L = 100;  
k = [1:1:L];  
cL = sum(1./k);  
pk = (1/cL)./k;  
figure;  
plot(k, discrete_pdf(k, k, pk));  
figure;  
plot(k, discrete_cdf(k, k, pk));  
figure;  
plot(k, 1-discrete_cdf(k, k, pk));
```

```
L = 1000;  
k = [1:1:L];  
cL = sum(1./k);  
pk = (1/cL)./k;  
figure;  
plot(k, discrete_pdf(k, k, pk));  
figure;  
plot(k, discrete_cdf(k, k, pk));  
figure;  
plot(k, 1-discrete_cdf(k, k, pk));
```

(b)

```
m = 20;  
k = [1:1:m];  
pk = (1/2).^k;  
figure;  
semilogy(k, pk);
```

3.80 The following Octave commands will plot the Lorenze curves:

```
L = 10;  
k = [1:1:L];  
cL = sum(1./k);  
pk = (1/cL)./k;  
Wk = k./L;  
Fk = discrete_cdf(k, k, pk);  
figure;  
plot(Fk, Wk);
```

```
L = 100;  
k = [1:1:L];  
cL = sum(1./k);  
pk = (1/cL)./k;  
Wk = k./L;  
Fk = discrete_cdf(k, k, pk);  
figure;  
plot(Fk, Wk);
```

```
L = 1000;  
k = [1:1:L];  
cL = sum(1./k);  
pk = (1/cL)./k;  
Wk = k./L;  
Fk = discrete_cdf(k, k, pk);  
figure;  
plot(Fk, Wk);
```

3.81 The following Octave commands will plot the requested curves:

```
figure;
hold on;
n = 100; p = 0.1;
k = [0:1:n];
Wk = zeros(1, n+1);
for i = 0:n,
    v = [0:i];
    Wk(i+1) = sum(v.*binomial_pdf(v, n, p))./(n*p);
end;
Fk = binomial_cdf(k, n, p);
plot(Fk, Wk);

n = 100; p = 0.5;
k = [0:1:n];
Wk = zeros(1, n+1);
for i = 0:n,
    v = [0:i];
    Wk(i+1) = sum(v.*binomial_pdf(v, n, p))./(n*p);
end;
Fk = binomial_cdf(k, n, p);
plot(Fk, Wk);

n = 100; p = 0.9;
k = [0:1:n];
Wk = zeros(1, n+1);
for i = 0:n,
    v = [0:i];
    Wk(i+1) = sum(v.*binomial_pdf(v, n, p))./(n*p);
end;
Fk = binomial_cdf(k, n, p);
plot(Fk, Wk);
```

3.82

- (a)
- (b)
- (c)

3.83 The following Octave commands will generate the requested samples of the Zipf random variable and the requested plots.

```
L = 10;  
k = [1:1:L];  
cL = sum(1./k);  
pk = (1/cL)./k;  
Sk = discrete_rnd(200, k, pk);  
figure;  
plot(Sk);  
figure;  
hist(Sk, k);
```

```
L = 100;  
k = [1:1:L];  
cL = sum(1./k);  
pk = (1/cL)./k;  
Sk = discrete_rnd(200, k, pk);  
figure;  
plot(Sk);  
figure;  
hist(Sk, k);
```

```
L = 1000;  
k = [1:1:L];  
cL = sum(1./k);  
pk = (1/cL)./k;  
Sk = discrete_rnd(200, k, pk);  
figure;  
plot(Sk);  
figure;  
hist(Sk, k);
```

3.84 The following Octave commands generate the samples of the St. Peter's Paradox random variable and the requested plots.

```
m = 20;  
k = [1:1:m];  
pk = (1/2).^k;  
Sk = discrete_rnd(200, k, pk);  
figure;  
plot(Sk);  
figure;  
hist(Sk, k);
```

3.85 The following Octave commands generate the requested pairs and plots:

(a)

```
k = [1:10];  
pk = ones(1,10)./10;  
Sx = discrete_rnd(200, k, pk);  
Sy = discrete_rnd(200, k, pk);  
figure;  
hist(Sx, k);  
figure;  
hist(Sy, k);
```

(b)

```
Sz = Sx + Sy;  
figure;  
hist(Sz, [2:20]);
```

(c)

```
Sw = Sx .* Sy;  
figure;  
hist(Sw, 10);
```

(d)

```
Sv = Sx ./ Sy;  
figure;  
hist(Sv, 10);
```

3.86 The following Octave commands generate the requested pairs and plots:

(a)

```
Sx = binomial_rnd(8, 0.5, 1, 200);  
Sy = binomial_rnd(4, 0.5, 1, 200);  
figure;  
hist(Sx, [0:8]);  
figure;  
hist(Sy, [0:4]);
```

(b)

```
Sz = Sx + Sy;  
figure;  
hist(Sz, [0:12]);
```


3.87 The following Octave commands generate the requested pairs and plots:

(a)

```
Sx = poisson_rnd(5, 1, 200);  
Sy = poisson_rnd(10, 1, 200);  
figure;  
hist(Sx, [0:15]);  
figure;  
hist(Sy, [0:20]);
```

(b)

```
Sz = Sx + Sy;  
figure;  
hist(Sz, [0:35]);
```

Problems Requiring Cumulative Knowledge

3.88

a)

$$P[\text{pass the test}] = (1 - p) + p(1 - \alpha)$$

$$P[\text{fail the test}] = p\alpha$$

$$P[k \text{ items}] = [(1 - p) + p(1 - \alpha)]^{k-1}(p\alpha)^1$$

b)

3.89

The number of transmissions is a geometric RV. The average number of transmissions is:

$$\begin{aligned} \sum_{k=1}^{\infty} k p^{k-1} (1 - p) &= (1 - p) \sum_{k=1}^{\infty} \frac{d p^k}{d p} \\ &= (1 - p) \frac{d}{d p} \sum_{k=1}^{\infty} p^k \\ &= (1 - p) \frac{d}{d p} \frac{1}{1 - p} \\ &= \frac{1}{1 - p} \end{aligned}$$

The message transmission takes $\frac{2T}{1 - P}$ seconds on the average. The maximum possible rate = $(1 - P)/2T$.

3.90) We want to find n so that the n th arrival is after more than 2 minutes 90% of the time:

$$P[N(2) \leq n] = 0.90 = \sum_{k=0}^n \frac{2^k}{k!} e^{-2}$$

By trial and error we find $n=5$.

3.91

$$\begin{aligned}
 & \text{58 a) } P[\text{signal present}|X = k] \\
 &= \frac{P[\text{signal present}, X = k]}{P[X = k|\text{signal present}]P[\text{present}] + P[X = k|\text{signal absent}]P[\text{absent}]} \\
 &= \frac{\frac{\lambda_1^k}{k!}e^{-\lambda_1}p}{\frac{\lambda_1^k}{k!}e^{-\lambda_1}p + \frac{\lambda_0^k}{k!}e^{-\lambda_0}(1-p)} \\
 &= \frac{\lambda_1^k e^{-\lambda_1} p}{\lambda_1^k e^{-\lambda_1} p + \lambda_0^k e^{-\lambda_0} (1-p)}
 \end{aligned}$$

Similarly,

$$P[\text{signal absent}|X = k] = \frac{\lambda_0^k e^{-\lambda_0} (1-p)}{\lambda_1^k e^{-\lambda_1} p + \lambda_0^k e^{-\lambda_0} (1-p)}$$

b) Decide signal present if $P[\text{signal present}|X=k] > P[\text{signal absent}|X=k]$, i.e.,

$$\begin{aligned}
 & \lambda_1^k e^{-\lambda_1} p > \lambda_0^k e^{-\lambda_0} (1-p) \\
 & \left(\frac{\lambda_1}{\lambda_0}\right)^k > \frac{1-p}{p} e^{\lambda_1 - \lambda_0} \quad (\lambda_1 > \lambda_0) \\
 & k > \frac{\ln \frac{1-p}{p} + \lambda_1 - \lambda_0}{\ln \lambda_1 - \ln \lambda_0}
 \end{aligned}$$

The threshold T is

$$T = \frac{\ln \frac{1-p}{p} + \lambda_1 - \lambda_0}{\ln \lambda_1 - \ln \lambda_0}$$

c)

$$\begin{aligned}
 P_e &= P[X < T|\text{signal present}]P[\text{present}] + P[X > T|\text{signal absent}]P[\text{absent}] \\
 &= p \sum_{k=0}^{[T]} \frac{e^{-\lambda_1} \lambda_1^k}{k!} + (1-p) \sum_{k=[T]}^{\infty} \frac{e^{-\lambda_0} \lambda_0^k}{k!}
 \end{aligned}$$

3.92

a)

$$\begin{aligned}
 P[\text{prefix has } k \text{ 0s}] &= P[kM \leq n \leq kM + M - 1] \\
 &= \sum_{kM}^{kM+M-1} p^n (1-p) \\
 &= (1-p)p^{kM} (1+p+\dots+p^{M-1}) \\
 &= p^{kM} (1-p^M) \\
 &= \left(\frac{1}{2}\right)^k \left(1 - \frac{1}{2}\right) \\
 &= \left(\frac{1}{2}\right)^{k+1}
 \end{aligned}$$

$$\begin{aligned}
 \text{b) } E[L] &= E[k] + 1 + m \\
 &= \sum_{k=0}^{\infty} k \left(\frac{1}{2}\right)^{k+1} + 1 + m \\
 &= m + 2
 \end{aligned}$$

$$\begin{aligned}
 \text{c) } E[\text{run length (including one 1 at the end)}] &= \sum_0^{\infty} (n+1)p^n(1-p) \\
 &= (1-p) \sum_0^{\infty} \frac{d}{dp} p^{n+1} \\
 &= (1-p) \frac{d}{dp} \sum_0^{\infty} p^{n+1} \\
 &= (1-p) \frac{d}{dp} \frac{p}{1-p} \\
 &= \frac{1}{1-p}
 \end{aligned}$$

$$\text{Compression ratio} = \frac{\frac{1}{1-p}}{m+2} = \frac{1}{(1-p)(m+2)}$$